Dealing with Disappearance of nodes in Social Networks

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October 2012

Outline

- Motivation
- Objectives
- Basic concepts
- Managing single node disappearance
- Managing group disappearance
- Validation
- Conclusion and future work



- May help to find
 - Relevant and personalized information by following/exploring the ties between elements
 - Groups and communities within a network
 - Influential actors in a group/network
 - Network evolution and prediction
 - etc.



- Knowing the groups and communities
 - Helps understand the shared interests, status, ...
 - what kind of products or services to suggest?
 - How to communicate with actors?
- Identifying the influential actors
 - Can enable to gauge their impact (or influence) on the other employees or clients
 - Predict the consequences of their departure
- Understanding the dynamics of the network structure
 - allows to forecast the behavior of clients and plan new changes and products

Motivation (I)

- The network structure is very dynamic
 - entities and links appear and disappear frequently
 - Ex: social networking sites, companies, school, ...
- Some nodes have more influence than others (i.e. central nodes)
 - they have the best ability to spread the information flow within the network
- The disappearance of a central node does not have the same impact as the deletion of a rather peripheral one
- Non-influential nodes, grouped as a set, may play an important role in the information flow or the network connectivity.

Objectives

- Predict the network evolution after a node/ group disappearance
 - Manage the network structure in such a way that a breakdown will never happen even if an actor set disappears
 - add new links
 - select a group/node substitute based on the leaving group configurations (or node role)

Related work

- Several studies deal with link prediction and/or network evolution in a social network
 - Liben-Nowell & Kleinberg (2003)
 - Exploit the proximity of nodes within a network to predict network structure by adding new edges.
 - Find a substitute to the deleted node provided it plays a role in the network. New links can then be established according to some predefined options (e.g. clique, ring).
 - Missaoui et *al.* (2011) Sarr et Missaoui (2012)
 - predict the new structure of a social network once a node disappears
 - possibly find a substitute for the leaving node provided it plays a role in the network and/or the information flow quality.
- Most of the proposed solutions
 - can explode the number of links within the network.
 - do not deal with a group disappearance that can happen in any social network

Basic concepts

Some definitions

- Degree centrality
 - provides the number of direct links of a node
 - it helps identify leaders which have the highest number of links within the network.
- Betweenness centrality
 - expresses the amount of control that a node (or a group of vertices) possesses over the interactions of other nodes in the network.
 - It is high for mediators (or brokers) which are nodes that act as intermediaries between other nodes or as joins between communities.
- Closeness centrality
 - indicates how a node is close to the other nodes in the network and hence how fast information circulates from a given node to other reachable nodes in the network.

Node Classes

- Nodes inside a group are classified into two categories based on their centrality measures (degree, betweenness, and closeness centrality)
 - Critical nodes
 - Nodes playing central roles: i.e. nodes with the highest centrality measures
 - Three types of critical nodes : *leader, mediator* and *witness*
 - critical disappearance : critical nodes that are unique in their class
 - Non critical nodes
 - Nodes playing less central roles
 - Two possible types
 - *Finger* : node whose centrality measure deviates slightly from the one of a critical node
 - Follower : neither critical nor finger node
 - non-critical disappearance : either non-critical deleted node or a critical node which is not unique in its class

Node Classes





If we set a deviation of 20%, then node 2, 3 and 9 will be the *fingers* of the leader node 4.

Information Flow Quality

- Hypothesis
 - Assume S a connected, undirected and unweighted network
- Eccentricity of a node *i* : *ε* (*i*)
 - Greatest geodesic distance between *i* and any other node *j* in the network
- Information flow degree of S : λ_F (S)
 - Total number *N* of nodes that can receive the information
- Information flow (IF) quality of $S : Y_F(S)$
 - $Y_F(S) = \langle \lambda_F(S), \varepsilon(w) \rangle$
 - ε (w) is the eccentricity of the witness that is the best entry of the network based on its closeness centrality

Dealing with single node disappearance

Managing a single node disappearance?

- Main idea
 - After the disappearance of a node, add new links and possibly find a substitute if the information flow quality is not ensured (e.g., disconnected network, delays in information propagation)



Managing a single node disappearance?

• Main idea

- After the disappearance of a node, add new links and possibly find a substitute if the information flow quality is not ensured
 - i.e. if the network is disconnected



Requirements

- Let S' be the network obtained after
 - (i) deleting node Ni from S,
 - (ii) removing links attached to Ni,
 - and (iii) adding new links.
- The idea is to minimize the number of added links while maintaining the following compound condition:

$$\frac{\lambda_F(S) - \lambda_F(S')}{\lambda_F(S)} \le \varepsilon_1 \quad \wedge \quad \frac{\epsilon_{S'}(w) - \epsilon_S(w)}{\epsilon_S(w)} \le \varepsilon_2$$

In other words, after a node disappearance, what is the "optimal" update that maintains the IF quality w.r.t some given deviations ε_1 and ε_2

Dealing with critical disappearance

- A critical disappearance of *Nc* is handled as follows:
 - Nc has some fingers (at least one). In this case, the first finger (F1) (finger with the highest centrality measure) is chosen as the substitute
 - 2. There is no finger to replace the leaving critical node. Then, one of the remaining critical nodes is selected as a substitute
 - The critical node with the shortest geodesic distance to Nc is used as a substitute
 - 3. There is no other critical nodes, which means that the leaving node was the only critical node (e.g., a central node of a star network). Then, one of the Nc neighbors with the highest centrality measure is selected as the substitute.
- For each one of the three cases, any neighbor of Nc is linked to the substitute if ever it has not been already linked

Critical disappearance – JOAN-C



Node 3 is the substitute since it may be the finger and/or the neighbor with the highest centrality measure

Dealing with non-critical disappearance

- A non-critical disappearance of *Ni* is handled as follows:
 - 1. The degree of *Ni* is equal to one, i.e., *Ni* has one neighbor.
 - Nothing to do because both conditions in the previous Formula are true for a network of size at least equal to $1/\epsilon_1$.
 - 2. The degree of Ni is greater than one.
 - While the tolerated information flow degree is not ensured, a link between two neighbors Nj and Nk of Ni is added as follows:
 - 1. A direct link is added if *Nj* or *Nk* is a critical node.
 - 2. An indirect link (between *Nj* and *Nk*) is added if neither *Nj* nor *Nk* is a critical one. The indirect link is set between *Nj* and *Nk* through another neighbor *Ns* which has the shortest path to a critical node *Nc*.



Non-critical disappearance - JOAN (2)



Managing a group disappearance

Configurations of a group

- A group G, can have one of the 3 configurations
 - Scattered
 - there is no direct link for any couple of nodes in G
 - Continuous group
 - there exists a direct link for any couple of nodes (Ni, Nk)
 - There exists an indirect link between Ni and Nk through only nodes inside G.
 - Hybrid group
 - Combination of both previous configurations coexist

Example







Scattered group

- each node N_i in G is far from the other ones
 - its disappearance will not involve another node N_k in the group.
 - i.e. by using JOAN or JOAN-C, we can face the disappearance of N_i without using any $N_k \subseteq G$ as a substitute or linking directly N_i to N_k .
- For each node $N_i \in G$,
 - apply the JOAN or JOAN-C algorithm based on the class of N_i to manage its disappearance.

Contiguous group

- disappearance of a contiguous group is handled in one shot.
 - JOAN or JOAN-C not efficient
 - may lead to tie two nodes that will leave afterwards because they belong to the same leaving group
- Every leaving group is replaced by another one to:
 - preserve the overall shape of the network
 - and ensure the information dissemination
- To replace a group G, we consider all neighbors Γ(G) of any node in G.
 - Let $\Gamma(Ni)$ be the set of neighbors of a node Ni

$$\Gamma(G) = \{\bigcup_i \Gamma(N_i) \setminus G \mid N_i \in G\}$$

Contiguous group : common neighbors

- To know the nodes to consider first when finding the substitute of a leaving group, we define common neighbors.
- Total common neighbors $\Gamma_{T}(G)$
 - the set of nodes that belong to each neighborhood of a node in G

 $\Gamma_T(G) = \{\bigcap_i \Gamma(N_i) \mid N_i \in \Gamma(G)\}$

- Partial common neighbors $\Gamma_P(G)$
 - The union of the intersections of the neighbors of any couple of nodes in G.

$$\Gamma_P(G) = \bigcup_{i,k} \{ \Gamma(N_i) \cap \Gamma(N_k) \mid (N_i, N_k) \in \Gamma(G) \}$$

Finding a substitute

- Let G' = substitute of a leaving group G
- Γ_T(G) ≠ Ø : there is at least one common neighbor for all nodes in G
 Γ_T(G) is used as a substitute group
- $\Gamma_{T}(G) = \emptyset \land \Gamma_{P}(G) \neq \emptyset$
 - Set G' by choosing *n* nodes in $\Gamma_{P}(G)$ in such a way that
 - |G| = |G'|
 - G' has the highest group degree centrality
- $\Gamma_{T}(G) = \emptyset \land \Gamma_{P}(G) = \emptyset$
 - Set G' by gathering a set of nodes from $\Gamma(G)$ in order to get a centrality value similar to or higher than the group degree centrality of G

Connecting G' to G neighborhood

- After finding G':
 - add direct links between any couple of nodes in G' if ever they are not yet directly or indirectly linked.
 - for any node $N_i \in \Gamma(G)$ and $N_i \notin G'$, we add a link between N_i and only one of the nodes in the substitute group G'.

Example



Other group configurations

- Hybrid
 - first identify the scattered and contiguous subgroups,
 - for each subgroup, we apply one of the procedures presented earlier.
- Group disappearance with identified communities
 - a leaving group can be localized either in one single community or spread over many communities
 - When the group is distributed over several communities, we first identify such communities and the corresponding subgroups.
 - For each subgroup within a community, we apply the procedure based on its nature (e.g., scattered)
 - When the overall group of the leaving node belongs to the same community, apply the procedure described above for a subgroup

Validation

- Goal
 - Evaluate the performance in terms of execution time and number of added links when node disappearance becomes more frequent and group configurations vary.

• Data sets

- Autonomous systems AS-733 of the Stanford University dataset (2009)
 - Undirected and unweighted network with 6474 nodes and 13233 edges.
 - describes a graph of Internet routers with a communication network model of who-talksto-whom

• Platform

- A prototype that uses NetworkX and the Python language.
- Intel Core i5 with 8 GB of RAM and 3.20 GHz running under Linux Ubuntu
- cloud infrastructure PiCloud

Single node disappearance results

- Overall performance
 - For a deviation = 0 and sparse graphs
 - Average added links : 5 links per disappearance
 - Average response time : less than 80 ms
- Impact of the information flow deviation
 - beyond a deviation of 2,5%
 - Average added links is less than 1 per disappearance
 - The eccentricity decreases at least by a factor of 60%
 - For more than 60% of deletions, no network update because the requirements are already ensured.





http://larim.uqo.ca/SNA/SarrSNAM.pdf

Group disappearance results

- Overall performance
 - Low increase of the number of added links
 - Average added links : 0,2 links per disappearance with 40% of disappearance.
 - Response time depends on the network density:
 - sparse graph (with hybrid group) leads to less added links and is less time consuming than dense graph.
 - Community identification gives more performances
 - Communities lead to parallelism and require less than 5ms where lack of communities consume 50ms to update network after a group disappearance.

Validation graphs



Future work

- Ongoing work:
 - we plan a more pro-active to notify and recommend ways to strengthen the weak links.
- Consider additional features such as the intensity of interactions among nodes and the direction of links to better capture the role of nodes in information dissemination
- Characterize influential actors in the network based on the roles we defined earlier (e.g., a witness and its top ranked fingers).